

# Heuristic Approach for Bicriteria in Constrained Three Stage Flow Shop Scheduling

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## Abstract

This paper presents bicriteria in n-jobs, three machines flow shop scheduling problem to minimize the total elapsed time and rental cost of the machines under a specified rental policy in which the processing time, independent setup time each associated with probabilities including transportation time and job block concept. Further the concept of the break down interval for which the machines are not available for the processing is included. A heuristic approach method to find optimal or near optimal sequence has been discussed. A computer programme followed by a numerical illustration is given to substantiate the algorithm.

**Keywords:** Flow Shop, Processing time, Setup time, Makespan, Break-down interval, Job block, Transportation time, Rental Cost.

## 1. Introduction

Scheduling is broadly defined as the process of the allocation of resources over time to perform a collection of tasks. Scheduling problems in their simple static and deterministic forms are extremely simple to describe and formulate, but are difficult to solve because they involve complex combinatorial optimization. For example, if n jobs are to be performed on m machines, there are potentially  $(n!)^m$  sequences, although many of these may be infeasible due to various constraints. Single criterion is deemed as insufficient for real and practical applications. Thus considering problems with more than one criterion is a practical direction of research for real-life scheduling problems. The bicriteria scheduling problems are motivated by the fact that they are more meaningful from practical point of view. The classical scheduling literature commonly assumes that the machines are never unavailable during the process. But there are feasible sequencing situations where machines while processing the jobs get sudden break-down due to failure of a component of machines for a certain interval of time or the machines are supposed to stop their working for a certain interval of time due to some external imposed policy such as stop of flow of electric current to the machines by a government policy due to shortage of electricity production. In each case this may be well observed that working of machines is not continuous and is subject to breakdown for certain interval of time. The majority of scheduling research assumes setup as negligible or part of processing time. While this assumption adversely affects solution quality for many applications which require explicit treatment of setup. Such applications have motivated increasing interest to include setup considerations in scheduling theory. One of the earliest results in flow shop scheduling theory is an algorithm given by Johnson [1954] for scheduling jobs in a two machine flowshop to minimize the time at which all jobs are completed. Smith [1967] considered minimisation of mean flow time and maximum tardiness. Some of the noteworthy heuristic approaches are due to Maggu & Das [1977], Yoshida & Hitomi [1979], Singh T.P. [1985], Adiri [1989], Akturk & Gorgulu [1999], Brucker and S.Knust [2004], Chandramouli [2005], N Chikhi [2008], Belwal and Mittal [2008], Khodadadi A. [2008], Pandian & Rajendran [2010] by considering various parameters. Gupta & Sharma [2011] studied bicriteria in n x 3 flow shop scheduling under specified rental policy, processing time associated with probabilities including transportation time and job block criteria. We have extended the study made by Gupta and Sharma [2011] by introducing the concept of setup time and breakdown interval. This paper considers a more practical scheduling situation in which certain ordering of jobs are prescribed either by technological constraints or by externally imposed policy.

## 2. Practical Situation

Many applied and experimental situations exist in our day-to-day working in factories and industrial production concerns etc. When the machines on which jobs are to be processed are planted at different places, the transportation time (which includes loading time, moving time and unloading time etc.) has a significant role in production concern. Setup includes work to prepare the machine, process or bench for product parts or the cycle. This includes obtaining tools, positioning work-in-process material, return tooling, cleaning up, setting the required jigs and fixtures, adjusting tools and inspecting material and hence significant. Various practical situations occur in real life when one has got the assignments but does not have one's own machine or does not have enough money or does not want to take risk of investing huge amount of money to purchase machine. Under such circumstances, the machine has to be taken on rent in order to complete the assignments. In his starting career, we find a medical practitioner does not buy expensive machines say X-ray machine, the Ultra Sound Machine, Rotating Triple Head Single Positron Emission Computed Tomography Scanner, Patient Monitoring Equipment, and Laboratory Equipment etc., but instead takes on rent. Rental of medical equipment is an affordable and quick solution for hospitals, nursing homes, physicians, which are presently constrained by the availability of limited funds due to the recent global economic recession. Renting enables saving working capital, gives option for having the equipment, and allows upgradation to new technology. Further the priority of one job over the other may be significant due to the relative importance of the jobs. It may be because of urgency or demand of that particular job. Hence, the job block criteria become important. Another event which is mostly considered in the models is the break-down of machines. There may also be delays due to material, changes in release and tail dates, tools unavailability, failure of electric current, the shift pattern of the facility and fluctuations in processing times. All of these events complicate the scheduling problem in most cases. Hence the criterion of break-down interval becomes significant.

### 3. Notations

- $S$  : Sequence of jobs 1,2,3,...,n
- $S_k$  : Sequence obtained by applying Johnson's procedure,  $k = 1, 2, 3, \dots$
- $M_j$  : Machine  $j, j = 1, 2, 3$
- $M$  : Minimum makespan
- $a_{ij}$  : Processing time of  $i^{\text{th}}$  job on machine  $M_j$
- $p_{ij}$  : Probability associated to the processing time  $a_{ij}$
- $s_{ij}$  : Set up time of  $i^{\text{th}}$  job on machine  $M_j$
- $q_{ij}$  : Probability associated to the set up time  $s_{ij}$
- $A_{ij}$  : Expected processing time of  $i^{\text{th}}$  job on machine  $M_j$
- $S_{ij}$  : Expected set up time of  $i^{\text{th}}$  job on machine  $M_j$
- $L$  : Length of the break-down interval
- $A'_{ij}$  : Expected processing time of  $i^{\text{th}}$  job after break-down effect on machine  $M_j$
- $\beta$  : Equivalent job for job – block
- $C_i$  : Rental cost of  $i^{\text{th}}$  machine
- $L_f(S_k)$ : The latest time when machine  $M_j$  is taken on rent for sequence  $S_k$
- $t_{ij}(S_k)$ : Completion time of  $i^{\text{th}}$  job of sequence  $S_k$  on machine  $M_j$
- $t'_{ij}(S_k)$  : Completion time of  $i^{\text{th}}$  job of sequence  $S_k$  on machine  $M_j$  when machine  $M_j$  start processing jobs at time  $L_f(S_k)$
- $T_{i,j \rightarrow k}$ : Transportation time of  $i^{\text{th}}$  job from  $j^{\text{th}}$  machine to  $k^{\text{th}}$  machine
- $I_{ij}(S_k)$ : Idle time of machine  $M_j$  for job  $i$  in the sequence  $S_k$
- $U_j(S_k)$ : Utilization time for which machine  $M_j$  is required, when  $M_j$  starts processing jobs at time  $L_f(S_k)$
- $R(S_k)$  : Total rental cost for the sequence  $S_k$  of all machine

#### 4. Rental Policy (P)

The machines will be taken on rent as and when they are required and are returned as and when they are no longer required i.e. the first machine will be taken on rent in the starting of the processing the jobs, 2<sup>nd</sup> machine will be taken on rent at time when 1<sup>st</sup> job is completed on 1<sup>st</sup> machine and transported to 2<sup>nd</sup> machine, 3<sup>rd</sup> machine will be taken on rent at time when 1<sup>st</sup> job is completed on the 2<sup>nd</sup> machine and transported.

#### 5. Problem Formulation

Let some job  $i$  ( $i = 1, 2, \dots, n$ ) are to be processed on three machines  $M_j$  ( $j = 1, 2, 3$ ) under the specified rental policy P. Let  $a_{ij}$  be the processing time of  $i^{\text{th}}$  job on  $j^{\text{th}}$  machine with probabilities  $p_{ij}$  and  $s_{ij}$  be the setup time of  $i^{\text{th}}$  job on  $j^{\text{th}}$  machine with probabilities  $q_{ij}$ . Let  $A_{ij}$  be the expected processing time and  $S_{i,j}$  be the expected setup time of  $i^{\text{th}}$  job on  $j^{\text{th}}$  machine. Let  $T_{i,j \rightarrow k}$  be the transportation time of  $i^{\text{th}}$  job from  $j^{\text{th}}$  machine to  $k^{\text{th}}$  machine. Our aim is to find the sequence  $\{S_k\}$  of the jobs which minimize the rental cost of all the three machines while minimizing total elapsed time.

The mathematical model of the problem is as shown in table 1.

Minimize  $U_j(S_k)$  and

Minimize  $R(S_k) = t_{n1}(S_k) \times C_1 + U_2(S_k) \times C_2 + U_3(S_k) \times C_3$

Subject to constraint: Rental Policy (P)

Our objective is to minimize rental cost of machines while minimizing total elapsed time.

#### 6. Algorithm

**Step 1:** Calculate the expected processing times and expected set up times as follows

$$A_{ij} = a_{ij} \times p_{ij} \quad \text{and} \quad S_{ij} = s_{ij} \times q_{ij} \quad \forall i, j = 1, 2, 3$$

**Step 2:** Check the condition

$$\text{Either } \text{Min}\{A_{i1} + T_{i,1 \rightarrow 2} - S_{i2}\} \geq \text{Max}\{A_{i2} + T_{i,1 \rightarrow 2} - S_{i1}\} \\ \text{or} \quad \text{Min}\{A_{i3} + T_{i,2 \rightarrow 3} - S_{i2}\} \geq \text{Max}\{A_{i2} + T_{i,2 \rightarrow 3} - S_{i3}\} \text{ or both for all } i$$

If the conditions are satisfied then go to step 3, else the data is not in the standard form.

**Step 3:** Introduce the two fictitious machines G and H with processing times  $G_i$  and  $H_i$  as

$$G_i = A_{i1} + A_{i2} + \max(S_{i1}, S_{i2}) + T_{i,1 \rightarrow 2} \quad \text{and} \quad H_i = A_{i2} + A_{i3} - S_{i3} + T_{i,2 \rightarrow 3}$$

**Step 4:** Find the expected processing time of job block  $\beta = (k, m)$  on fictitious machines G & H using equivalent job block criterion given by Maggu & Das [1977]. Find  $G_\beta$  and  $H_\beta$  using

$$G_\beta = G_k + G_m - \min(G_m, H_k) \quad \text{and} \quad H_\beta = H_k + H_m - \min(G_m, H_k)$$

**Step 5:** Define new reduced problem with processing time  $G_i$  &  $H_i$  as defined in step 3 and replace job block  $(k, m)$  by a single equivalent job  $\beta$  with processing times  $G_\beta$  &  $H_\beta$  as defined in step 4.

**Step 6:** Using Johnson's procedure, obtain all sequences  $S_k$  having minimum elapsed time. Let these be  $S_1, S_2, \dots, S_r$

**Step 7:** Prepare In – Out tables for the sequences obtained in step 6 and read the effect of break-down interval (a, b) on different jobs on the lines of Singh T.P. [1985].

**Step 8:** Form a reduced problem with processing times  $A'_{ij}$  ( $j=1, 2, 3$ )

If the break-down interval (a, b) has effect on job  $i$  then  $A'_{ij} = A_{ij} + L \quad \forall i, j = 1, 2, 3$

Where  $L = b - a$ , the length of break-down interval

If the break-down interval (a, b) has no effect on  $i^{\text{th}}$  job then  $A'_{ij} = A_{ij} \quad \forall i, j = 1, 2, 3$

**Step 9:** Now repeat the procedure to get optimal sequence  $S'_k$

**Step 10:** Prepare In – Out tables for  $S'_k$  and compute total elapsed time  $t_{n3}(S'_k)$

**Step 11:** Compute latest time  $L_3$  for machine  $M_3$  for sequence  $S'_k$  as

$$L_3(S'_k) = t_{n3}(S'_k) - \sum_{i=1}^n A'_{i3} - \sum_{i=1}^{n-1} S_{i,3}(S'_k)$$

**Step 12:** For the sequence  $S'_k$  ( $k = 1, 2, \dots, r$ ), compute

- I.  $t_{n2}(S'_k)$
- II.  $Y_1(S'_k) = L_3(S'_k) - A'_{1,2}(S'_k) - T_{1,2 \rightarrow 3}$
- III.  $Y_q(S'_k) = L_3(S'_k) - \sum_{i=1}^q A'_{i2}(S'_k) - \sum_{i=1}^q T_{i,2 \rightarrow 3} - \sum_{i=1}^{q-1} S_{i,2}(S'_k) + \sum_{i=1}^{q-1} A'_{i,3} + \sum_{i=1}^{q-1} T_{i,1 \rightarrow 2} + \sum_{i=1}^{q-2} S_{i3}(S'_k); q = 2, 3, \dots, n$
- IV.  $L_2(S'_k) = \min_{1 \leq q \leq n} \{Y_q(S'_k)\}$
- V.  $U_2(S'_k) = t_{n2}(S'_k) - L_2(S'_k)$ .

**Step 13:** Find  $\min \{U_2(S'_k)\}; k = 1, 2, \dots, r$

Let it be for the sequence  $S'_p$  and then sequence  $S'_p$  will be the optimal sequence.

**Step 14:** Compute total rental cost of all the three machines for sequence  $S'_p$  as:

$$R(S'_p) = t_{n1}(S'_p) \times C_1 + U_2(S'_p) \times C_2 + U_3(S'_p) \times C_3$$

## 7. Numerical Illustration

Consider 5 jobs, 3 machine flow shop problem with processing time, setup time associated with their respective probabilities and transportation time as given in table 2 and jobs 2 and 4 are processed as a group job (2, 4) with breakdown interval (12,14). The rental cost per unit time for machines  $M_1$ ,  $M_2$  and  $M_3$  are 2 units, 10 units and 8 units respectively, under the specified rental policy P. Our objective is to obtain an optimal schedule for above said problem to minimize the total production time / total elapsed time subject to minimization of the total rental cost of the machines.

**Solution: As per Step 1:** the expected processing times and expected setup times for machines  $M_1$ ,  $M_2$  and  $M_3$  are as shown in table 3.

**As per step 2 :** Here  $\text{Min} \{A_{i1} + T_{i,1 \rightarrow 2} - S_{i2}\} \geq \text{Max} \{A_{i2} + T_{i,1 \rightarrow 2} - S_{i1}\}$

**As per step 3:** The expected processing time for two fictitious machine G & H is as shown in table 4.

**As per Step 4:** Here  $\beta = (2, 4)$

$$G_\beta = 11.4 + 9.3 - 9.3 = 11.4$$

$$H_\beta = 9.8 + 6.8 - 9.3 = 7.3$$

**As per Step 5:** The reduced problem is as shown in table5.

**As per Step 6:** Using Johnson's method, the optimal sequence is

$$S = 3 - 5 - \beta - 1, \text{ i.e. } S = 3 - 5 - 2 - 4 - 1$$

**As per step 8:** The new processing times after breakdown effect are as shown in table 7

**As per step 9 :** Using Johnson's method optimal sequence is  $S'$

$$S' = 3 - 5 - 2 - 4 - 1$$

**As per step 10:** The In-Out table for the sequence  $S'$  is as shown in table 8.

Total elapsed time  $t_{n,3}(S') = 39.6$  units

**As per Step 11:**  $L_3(S') = t_{n3}(S') - \sum_{i=3}^n A'_{i,3} - \sum_{i=3}^{n-1} S_{i,3}(S')$   
 $= 39.6 - 15.8 - 1.4 - 2.5 = 9.8$  units

**As per Step 12:** For sequence  $S_2 = 1, 5, 9, 2, 6, 8 = 13.4$   
 $t_{n2}(S_2) = 33.8 - 14.8 + 16.4 = 17.4$   
 $Y_4 = 15.8 - 20.5 + 23.3 = 18.6$

The new reduced Bi-objective in 5. Out 25.12.2012 as in 10.

The latest possible time at which machine  $M_2$  should be taken on rent  $= L_2(S') = 9.8$  units.

Also, utilization time of machine  $M_2 = t_{n2}(S') - E_2(4)$  units  $17.4 - 9.8 = 7.6$

Total minimum rental cost  $= R(S') = t_{n1}(S') \times C_1 + U_2(S') \times C_2 + U_3(S') \times C_3$   
 $= 29.7 \times 2 + 7.6 \times 10 + 23.8 \times 8 = 489.8$  units

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**Tables**

Table 1 : The mathematical of the given problem

Jobs	Machine A				$T_{i,1 \rightarrow 2}$	Machine B				$T_{i,2 \rightarrow 3}$	Machine C			
i	$a_{i1}$	$p_{i1}$	$s_{i1}$	$q_{i1}$		$a_{i2}$	$p_{i2}$	$s_{i2}$	$q_{i2}$		$a_{i3}$	$p_{i3}$	$s_{i3}$	$q_{i3}$
1	$a_{11}$	$p_{11}$	$s_{11}$	$q_{11}$	$T_{1,1 \rightarrow 2}$	$a_{12}$	$p_{12}$	$s_{12}$	$q_{12}$	$T_{1,2 \rightarrow 3}$	$a_{13}$	$p_{13}$	$s_{13}$	$q_{13}$
2	$a_{21}$	$p_{21}$	$s_{21}$	$q_{21}$	$T_{2,1 \rightarrow 2}$	$a_{22}$	$p_{22}$	$s_{22}$	$q_{22}$	$T_{2,2 \rightarrow 3}$	$a_{23}$	$p_{23}$	$s_{23}$	$q_{23}$
3	$a_{31}$	$p_{31}$	$s_{31}$	$q_{31}$	$T_{3,1 \rightarrow 2}$	$a_{32}$	$p_{32}$	$s_{32}$	$q_{32}$	$T_{3,2 \rightarrow 3}$	$a_{33}$	$p_{33}$	$s_{33}$	$q_{33}$
4	$a_{41}$	$p_{41}$	$s_{41}$	$q_{41}$	$T_{4,1 \rightarrow 2}$	$a_{42}$	$p_{42}$	$s_{42}$	$q_{42}$	$T_{4,2 \rightarrow 3}$	$a_{43}$	$p_{43}$	$s_{43}$	$q_{43}$
-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
n	$a_{n1}$	$p_{n1}$	$s_{n1}$	$q_{n1}$	$T_{n,1 \rightarrow 2}$	$a_{n2}$	$p_{n2}$	$s_{n2}$	$q_{n2}$	$T_{n,2 \rightarrow 3}$	$a_{n3}$	$p_{n3}$	$s_{n3}$	$q_{n3}$

Table 2: 5 jobs, 3 machine flow shop problem with processing time

Jobs	Machine $M_1$				$T_{i,1 \rightarrow 2}$	Machine $M_2$				$T_{i,2 \rightarrow 3}$	Machine $M_3$			
i	$a_{i1}$	$p_{i1}$	$s_{i1}$	$q_{i1}$		$a_{i2}$	$p_{i2}$	$s_{i2}$	$q_{i2}$		$a_{i3}$	$p_{i3}$	$s_{i3}$	$q_{i3}$
1	27	0.2	3	0.3	2	7	0.3	3	0.2	2	19	0.2	4	0.2
2	30	0.2	2	0.1	1	20	0.2	2	0.2	1	18	0.3	3	0.2
3	41	0.1	2	0.3	2	20	0.2	1	0.2	2	14	0.2	2	0.3
4	23	0.2	4	0.1	2	23	0.1	2	0.2	3	23	0.1	4	0.2
5	20	0.3	2	0.2	4	10	0.2	3	0.2	1	25	0.2	5	0.1

Table 3: The expected processing times and expected setup times

Jobs	$A_{i1}$	$S_{i1}$	$T_{i,1 \rightarrow 2}$	$A_{i2}$	$S_{i2}$	$T_{i,2 \rightarrow 3}$	$A_{i3}$	$S_{i3}$
1	5.4	0.9	2	2.1	0.6	2	3.8	0.8
2	6.0	0.2	1	4.0	0.4	1	5.4	0.6
3	4.1	0.6	2	4.0	0.2	2	2.8	0.6
4	4.6	0.4	2	2.3	0.4	3	2.3	0.8
5	6.0	0.4	4	2.0	0.6	1	5.0	0.5

Table 4: The expected processing time for two fictitious machine G & H

Jobs	$G_i$	$H_i$
1	10.4	7.1
2	11.4	9.8
3	10.7	8.2
4	9.3	6.8
5	12.6	7.5

Table 5: The reduced problem is

Jobs	$G_i$	$H_i$
1	10.4	7.1
$\beta$	11.4	7.3
3	10.7	8.2
5	12.6	7.5

Table 6: The In – Out table for the optimal sequence S is

Jobs	Machine $M_1$	$T_{i,1 \rightarrow 2}$	Machine $M_2$	$T_{i,2 \rightarrow 3}$	Machine $M_3$
i	In – Out		In – Out		In - Out
3	0 – 4.1	2	6.1 – 10.1	2	12.1 – 14.9
5	4.7 – 10.7	4	14.7 – 16.7	1	17.7 – 22.7
2	11.1 – 17.1	1	18.1 – 22.1	1	23.2 – 28.6
4	17.3 – 21.9	2	23.9 – 26.2	3	29.2 – 31.5
1	22.3 – 27.7	2	29.7 – 31.8	2	33.8 – 37.6

Table 7: The new processing times after breakdown effect

Jobs	$A'_{i1}$	$S_{i1}$	$T_{i,1 \rightarrow 2}$	$A'_{i2}$	$S_{i2}$	$T_{i,2 \rightarrow 3}$	$A'_{i3}$	$S_{i3}$
1	5.4	0.9	2	2.1	0.6	2	3.8	0.8
2	8.0	0.2	1	4.0	0.4	1	5.4	0.6
3	4.1	0.6	2	4.0	0.2	2	4.8	0.6
4	4.6	0.4	2	2.3	0.4	3	2.3	0.8
5	6.0	0.4	4	2.0	0.6	1	5.0	0.5

Table 8: The In-Out table for the sequence  $S'$  is

Jobs	Machine $M_1$	$T_{i,1 \rightarrow 2}$	Machine $M_2$	$T_{i,2 \rightarrow 3}$	Machine $M_3$
i	In – Out		In – Out		In - Out
3	0 – 4.1	2	6.1 – 10.1	2	12.1 – 16.9
5	4.7 – 10.7	4	14.7 – 16.7	1	17.7 – 22.7
2	11.1 – 19.1	1	20.1 – 24.1	1	25.1 – 30.5
4	19.3 – 23.9	2	25.9 – 28.2	3	31.2 – 33.5
1	24.3 – 29.7	2	31.7 – 33.8	2	35.8 – 39.6

Table 9: The new reduced Bi-objective In – Out table is

<b>Jobs</b>	Machine M <sub>1</sub>	$T_{i,1 \rightarrow 2}$	Machine M <sub>2</sub>	$T_{i,2 \rightarrow 3}$	Machine M <sub>3</sub>
i	In – Out		In – Out		In - Out
3	0 – 4.1	2	9.8 – 13.8	2	15.8 – 20.6
5	4.7 – 10.7	4	14.7 – 16.7	1	21.2 – 26.2
2	11.1 – 19.1	1	20.1 – 24.1	1	26.7 – 32.1
4	19.3 – 23.9	2	25.9 – 28.2	3	32.7 – 35.0
1	24.3 – 29.7	2	31.7 – 33.8	2	35.8 – 39.6



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