# AN EFFECTIVE APPROACH OF FOUR-STEP METHOD FOR OPTIMAL SOLUTION OF TRANSPORTATION PROBLEM 

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#### Abstract

Transportation problem (TP) in operation research is one of the most in use optimization technique to deal the problems that are related with transportation of goods from sources to destinations. Initial Basic Feasible Solution (IBFS) plays a vital role in TP which offers a way to obtain the optimal solution. The objective is to prevail the total transportation cost equivalent or nearer to optimal solution. In this paper, an effective approach of Four Step Method (FSM) for optimal solution of TP has been brought up in order to get optimal solution of TPs. In this method we construct the Maximum Column Table (MCT) and Maximum Row Table (MRT). Several problems has been solved using this method to get the optimal solution. The outcomes of proposed method are contrasted with results of North West Corner Method (NWCM), Least Cost Method (LCM) and Vogal's Approximation Method (VAM). It is observed that the proposed method is not only achieving better results but also overcoming the limitation of VAM.


Key Words: Transportation problem, Initial Basic Feasible Solution, Optimal solution
Linear programming problem

## 1. Introduction

Operation Research (OR) is a subject that uses application of progressed expository strategies to assist make superior decisions. The subject Operation Research was introduced during second World War and utilized for military techniques. During second World War a group of scientists from mathematics, physical and social sciences were grant to consider about of different army operations. This team was exceptionally successful in reducing cost of transporting the military equipment in different areas of war.

The transportation problem is the particular foreshadow of Linear Programming Problem (LPP) that was first presented by Hitchcock [1] in order to distribute the production of different numerous localities. After few years T.C. Koopmans [2] gave a method called 'Optimum Utilization of the Transportation System'. These two contributions are known as basement for the progress of transportation problem. The main ambition of Transportation problem was to minimize the cost of transportation of goods from $n$ origin to $m$ destinations.

Later, this type of problem was initially worked by Dantzig in 1951 and flourished by Charnes and Copper in 1953.Many scholars have worked on such type of problems. In 2014,
M.Tanvir [3] and U.k. Das [4] proposed different methods which give almost optimal solution. In 2015, S.Akpan [5] modified the VAM by using the idea of Standard Deviation of rows and columns on balanced TP which showed better results. In 2016, M.S.Uddin [6] made modification on LCM and M.M. Ahmed [7] proposed a method called Incessant Allocation Method for solving transportation problem, this method was then checked on both profit maximizing and cost minimizing problems and concluded that it is giving much better results. In 2017, M.J.uddin [8] developed an algorithm named as Weighted Cost Opportunity which showed better results as compared to regular LCM.In the same year, M.R.S. Shaikh [9] proposed an algorithm and M.B. Hossain [10] gave a new method by considering average penalty of rows and average penalty of columns. The results obtained from this method is more acceptable as compared to the VAM. In 2018, R.Kumar [11] proposed a new approach called Direct Sum Method (DSM) which gave better IBFS as compared to well-known methods. In 2019, many methods proposed but among them BCE [12] and PGM [13] were considered to give best results. In 2020, A.M.Khoso [14] made modification on LCM and Hussain [15] made modification on VAM. Both modifications are considered to give better results in less time.

There are three well known methods that are being used to get the Initial Basic feasible Solution (IBFS) named as North West Corner Method (NWCM), Least corner Method (LCM) and Vogal's Approximation Method (VAM). Among these three methods, VAM is considered to give best IBFS.

The three fundamental and well known methods to solve transportation problems are:

### 1.1 North West Corner Method (NWCM)

This is the basic and primary method that was used to minimize the total transportation cost. In this method we first allocate the northwest cell the required supply or demand and cross out that respective row or column. This process continues till the total cost is minimized

### 1.2 Least Cost Method (LCM)

This method is more effective than NWCM. In this method we allocate the lowest cost in transportation table the required supply or demand and cross out that respective row or column. This process continues till the total cost is minimized.

### 1.3. Vogal's Approximation Method (VAM)

This method is considered to give the best results among above methods. In this method we calculate penalty by subtracting the two consecutive lowest costs of each row and column. We allocate largest penalty the required supply or demand and cross out that respective satisfied row or column. This process continues till the total cost is minimized.

### 1.4. Proposed Four-Step Method

In this proposed method we calculate the Maximum Row Table (MRT) and Maximum Column table (MCT) by subtracting the largest cost of each row (Column) from respective costs of that row (Column). After that we add MRT and MCT and name this table as revised table. We allocate the largest cost of revised table and cross out that respective row or column. This process continues till all supplies and demands are satisfied. Then we make same positioned allotment in original table as revised table

Let $\mathrm{x}_{\mathrm{ij}} \geq 0$ be the amount transported from the source $i$ to the destination $j$, The mathematical formulation of the problem is stated as:

| Minimize | $\mathbf{Z}=\sum_{i=1}^{m} \sum_{j=1}^{n} C i j x i j$ | (Total transportation cost) |
| :--- | :--- | :--- |
| Subject to | $\sum_{j=1}^{n} x \mathrm{ij}=\mathrm{ai}$ | (Supply from sources) |
|  | $\sum_{i=1}^{m} x \mathrm{ij}=\mathrm{bj}$ | (Demand from destinations) |
|  | $\mathrm{x}_{\mathrm{ij}} \geq 0$, | for all $i$ and $j$ |

Where $\quad \mathrm{Z}$ : Total transportation cost to be minimized.
$\mathrm{C}_{\mathrm{ij}}$ : transportation cost of the goods from each sources i to destination j .
$\mathrm{x}_{\mathrm{ij}}$ : Amount of goods driven away from source i to destination j .
$a_{i}$ : Amount of supply at each source $i$.
$b_{j}$ : Amount of demand at each destination $j$.

Optimal solution of Transportation Model (TM) has two pace: initially we obtain an initial basic feasible solution (IBFS) using $\mathrm{m}+\mathrm{n}-1$ destinations, where n and m are origins and destinations respectively.

## 2. Methodology

In this section we will explain step by step process of proposed method.
Step 1: Calculate the Maximum Row (MRT ) and Column Table (MCT ) by subtracting the largest cost of each row (Column) from respective costs of that row (Column ). i.e. if $\mathrm{C}_{11}$ is largest cost in first row ( Column ) then

For MRT,

$$
\begin{aligned}
& \mathrm{C}_{11}{ }^{+}=\mathrm{C}_{11}-\mathrm{C}_{11} \\
& \mathrm{C}_{12}{ }^{+}=\mathrm{C}_{11}-\mathrm{C}_{12} \\
& \mathrm{C}_{13^{+}}=\mathrm{C}_{11}-\mathrm{C}_{13}
\end{aligned}
$$

## For MCT

$$
\begin{aligned}
& \mathrm{C}_{11^{-}}=\mathrm{C}_{11}-\mathrm{C}_{11} \\
& \mathrm{C}_{21^{-}}=\mathrm{C}_{11}-\mathrm{C}_{21} \\
& \mathrm{C}_{31^{-}}=\mathrm{C}_{11}-\mathrm{C}_{31}
\end{aligned}
$$

And so on

Table 2.1: Maximum Row Table ( MRT )

| Destination <br> Sources | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\cdots$ | $\mathrm{D}_{\mathrm{n}}$ | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{S}_{1}$ | $\mathrm{C}_{11^{+}}=\mathrm{C}_{12}-\mathrm{C}_{11}$ | $\mathrm{C}_{12^{+}}=\mathrm{C}_{12}-\mathrm{C}_{12}$ | $\mathrm{C}_{13^{+}}=\mathrm{C}_{12}-\mathrm{C}_{13}$ | $\cdots$ | $\mathrm{C}_{1 \mathrm{n}^{+}}=\mathrm{C}_{12}-\mathrm{C}_{1 \mathrm{n}}$ | $\mathrm{S}_{1}$ |
| $\mathrm{~S}_{2}$ | $\mathrm{C}_{21^{+}}=\mathrm{C}_{22}-\mathrm{C}_{21}$ | $\mathrm{C}_{22^{+}}=\mathrm{C}_{22}-\mathrm{C}_{22}$ | $\mathrm{C}_{23^{+}}=\mathrm{C}_{22}-\mathrm{C}_{23}$ | $\cdots$ | $\mathrm{C}_{2 \mathrm{n}^{+}}=\mathrm{C}_{22}-\mathrm{C}_{2 \mathrm{n}}$ | $\mathrm{S}_{2}$ |
| $\mathrm{~S}_{3}$ | $\mathrm{C}_{31^{+}}=\mathrm{C}_{33}-\mathrm{C}_{31}$ | $\mathrm{C}_{32^{+}}=\mathrm{C}_{33}-\mathrm{C}_{32}$ | $\mathrm{C}_{33^{+}}=\mathrm{C}_{33}-\mathrm{C}_{33}$ | $\cdots$ | $\mathrm{C}_{3 \mathrm{n}}{ }^{+}=\mathrm{C}_{33}-\mathrm{C}_{3 \mathrm{n}}$ | $\mathrm{S}_{3}$ |
|  | $\cdots$ | $\cdots$ | $\cdots$ | $\cdots$ | $\cdots$ | $\cdots$ |
| $\mathrm{~S}_{\mathrm{m}}$ | $\mathrm{C}_{\mathrm{m} 1^{+}=\mathrm{C}_{\mathrm{m} 2}-\mathrm{C}_{\mathrm{m} 1}}$ | $\mathrm{C}_{\mathrm{m} 2^{+}}=\mathrm{C}_{\mathrm{m} 2}-\mathrm{C}_{\mathrm{m} 2}$ | $\mathrm{C}_{\mathrm{m}^{+}}=\mathrm{C}_{\mathrm{m} 2}-\mathrm{C}_{\mathrm{m} 3}$ | $\cdots$ | $\mathrm{C}_{\mathrm{mn}}{ }^{+}=\mathrm{C}_{\mathrm{m} 2}-\mathrm{C}_{\mathrm{mn}}$ | $\mathrm{S}_{\mathrm{m}}$ |
| Demand | $\mathrm{q}_{1}$ | $\mathrm{q}_{2}$ | $\mathrm{q}_{3}$ | $\cdots$ | $\mathrm{q}_{\mathrm{n}}$ |  |

Table 2.2: Maximum Column Table ( MCT )

| Destination | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | ... | $\mathrm{D}_{\mathrm{n}}$ | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{S}_{1}$ | $\mathrm{C}_{11^{-}}=\mathrm{C}_{31}-\mathrm{C}_{11}$ | $\mathrm{C}_{12}{ }^{-}=\mathrm{C}_{42}-\mathrm{C}_{11}$ | $\mathrm{C}_{13}{ }^{\text {a }}=\mathrm{C}_{13}-\mathrm{C}_{13}$ | ... | $\mathrm{C}_{1 \mathrm{n}}=\mathrm{C}_{2 \mathrm{n}}-\mathrm{C}_{1 \mathrm{n}}$ | $\mathrm{S}_{1}$ |
| $\mathrm{S}_{2}$ | $\mathrm{C}_{21}{ }^{-}=\mathrm{C}_{31}-\mathrm{C}_{21}$ | $\mathrm{C}_{22}=\mathrm{C}_{42}-\mathrm{C}_{22}$ | $\mathrm{C}_{23}{ }^{\text {a }}=\mathrm{C}_{13}-\mathrm{C}_{23}$ | $\ldots$ | $\mathrm{C}_{2 \mathrm{n}^{-}}=\mathrm{C}_{2 \mathrm{n}}-\mathrm{C}_{2 \mathrm{n}}$ | $\mathrm{S}_{2}$ |
| $\mathrm{S}_{3}$ | $\mathrm{C}_{31}{ }^{-} \mathrm{C}_{31}-\mathrm{C}_{31}$ | $\mathrm{C}_{32}{ }^{-}=\mathrm{C}_{42}-\mathrm{C}_{32}$ | $\mathrm{C}_{33^{\circ}}=\mathrm{C}_{13}-\mathrm{C}_{33}$ | $\ldots$ | $\mathrm{C}_{3 \mathrm{n}}{ }^{\prime}=\mathrm{C}_{2 \mathrm{n}}-\mathrm{C}_{3 \mathrm{n}}$ | $\mathrm{S}_{3}$ |
|  | ... | ... | ... | $\ldots$ | $\ldots$ | $\ldots$ |
| $\mathrm{S}_{\mathrm{m}}$ | $\mathrm{C}_{\mathrm{ml}}{ }^{-}=\mathrm{C}_{31}-\mathrm{C}_{\mathrm{ml}}$ | $\mathrm{Cm}_{2}{ }^{-}=\mathrm{C}_{42}-\mathrm{C}_{\mathrm{m} 2}$ | $\mathrm{Cm}_{\mathbf{3} 3}=\mathrm{C}_{13}-\mathrm{C}_{\mathrm{m} 3}$ | ... | $\mathrm{Cmm}^{\text {n }}=\mathrm{C}_{2 \mathrm{n}}-\mathrm{C}_{\mathrm{mn}}$ | $\mathrm{S}_{\mathrm{m}}$ |
| Demand | $\mathrm{q}_{1}$ | $\mathrm{q}_{2}$ | $\mathrm{q}_{3}$ | $\ldots$ | $\mathrm{q}_{\mathrm{n}}$ |  |

Step 2: Add MRT and MCT and name the new table as revised table.

Table C: Revised Table $=$ MRT + MCT

| Destination | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | $\ldots$ | $\mathrm{D}_{\mathrm{n}}$ | Supply |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{S}_{1}$ | $\mathrm{C}_{11}{ }^{*}=\mathrm{C}_{11^{+}}+\mathrm{Clin}^{-}$ | $\mathrm{C}_{12}{ }^{*}=\mathrm{C}_{12}{ }^{+}+\mathrm{Cl}_{12}{ }^{-}$ | $\mathrm{C}_{13}{ }^{*}=\mathrm{C}_{13}{ }^{+}+\mathrm{C}_{13}{ }^{\text {a }}$ | $\ldots$ | $\mathrm{C}_{1 \mathrm{n}}{ }^{*}=\mathrm{C}_{1 \mathrm{n}}{ }^{+}+\mathrm{Cln}^{-}$ | $\mathrm{S}_{1}$ |
| $\mathrm{S}_{2}$ | $\mathrm{C}_{21}{ }^{*}=\mathrm{C}_{21}{ }^{+}+\mathrm{C}_{21}{ }^{-}$ | $\mathrm{C}_{22}{ }^{*}=\mathrm{C}_{22}{ }^{+}+\mathrm{C}_{22}{ }^{-}$ | $\mathrm{C}_{23}{ }^{*}=\mathrm{C}_{23}++\mathrm{C}_{23}{ }^{\text {a }}$ | $\ldots$ | $\mathrm{C}_{2 \mathrm{n}}{ }^{*}=\mathrm{C}_{2 \mathrm{n}}{ }^{+}+\mathrm{C}_{2 \mathrm{n}}{ }^{-}$ | $\mathrm{S}_{2}$ |
| $\mathrm{S}_{3}$ | $\mathrm{C}_{31}{ }^{*}=\mathrm{C}_{31}{ }^{+}+\mathrm{C}_{31}{ }^{-}$ | $\mathrm{C}_{32}{ }^{*}=\mathrm{C}_{32}{ }^{+}+\mathrm{C}_{32}{ }^{-}$ | $\mathrm{C}_{33}{ }^{*}=\mathrm{C}_{33}{ }^{+}+\mathrm{C}_{33}{ }^{-}$ | $\ldots$ | $\mathrm{C}_{3 \mathrm{n}}{ }^{*}=\mathrm{C}_{3 \mathrm{n}}{ }^{+}+\mathrm{C}_{3 \mathrm{n}}{ }^{\circ}$ | $\mathrm{S}_{3}$ |
|  | $\cdots$ | $\cdots$ | $\cdots$ | $\ldots$ | $\ldots$ | $\ldots$ |
| $\mathrm{S}_{\mathrm{m}}$ | $\mathbf{C o m}_{\mathbf{m}}{ }^{*}=\mathbf{C m 1}^{+}+\mathrm{C}_{\mathrm{m} 1}{ }^{-}$ | $\mathrm{Cm}_{2}{ }^{*}=\mathrm{Cm}_{2}{ }^{+}+\mathrm{C}_{\mathrm{m} 2}{ }^{-}$ | $\mathrm{C}_{\mathrm{m} 3}{ }^{*}=\mathrm{Cm}_{3}{ }^{+}+\mathrm{C}_{\mathrm{m} 3}{ }^{-}$ | $\cdots$ | $\mathrm{Cmn}^{*}=\mathrm{Cmn}^{+}+\mathrm{Cmn}^{-}$ | $\mathrm{S}_{\mathrm{m}}$ |
| Demand | $\mathrm{q}_{1}$ | $\mathrm{q}_{2}$ | $\mathrm{q}_{3}$ | $\ldots$ | $\mathrm{q}_{\mathrm{n}}$ |  |

Step 3: Allocate largest cost $\mathrm{U}_{\mathrm{ij}}$ (as shown in eq (1))of revised table a maximum supply $\mathrm{a}_{\mathrm{i}}$ (or demand $b_{j}$ ) and cross out the satisfied row of supply (or column of demand). Then in same table select second largest cost $U_{k l}$ of revised table and repeat the process until all supply $a_{i}$ and demand $b_{j}$ are satisfied. If there are similar costs in the revised table, then select the cost $\mathbf{P}$ which gives us maximum sum of supply and demand of respective cost as shown in eq 2.
$\mathrm{U}_{\mathrm{ij}}=\max \left(\mathrm{c}_{\mathrm{ij}}\right), \mathrm{i}, \mathrm{j}=1,2,3 \ldots$

$$
\begin{equation*}
\mathbf{P}=\max \left(\mathrm{a}_{\mathrm{i}}+\mathrm{b}_{\mathrm{j}}\right)- \tag{1}
\end{equation*}
$$

Step 4: Make the same positioned allotment in the original table as revised table.

$$
\operatorname{Min} Z=c_{11} x_{11}+c_{12} X_{12}+c_{13} X_{13}+c_{21} X_{21}+\cdots+c_{i j} X_{i j}
$$

## 3. Numerical Illustration

In this paper, we have examined a balanced transportation problems chosen from literature. The result of these problems are then compared with well-known methods named as NWCM , LCM and VAM. We solve example 1 step-by-step.
3.1. Example 1: Consider the following balanced transportation problem.

Table 3.1:Transportation Problem

|  | D1 | D2 | D3 | Sources |
| :---: | :---: | :---: | :---: | :---: |
| S1 | 6 | 8 | 4 | 14 |
| S2 | 4 | 9 | 8 | 12 |
| S3 | 1 | 2 | 6 | 5 |
| Destinations | 6 | 10 | 15 | 31 |

Step 1: For MRT, subtract 8, 9, 6 from each elements of $1^{\text {st }}, 2^{\text {nd }}$ and $3^{\text {rd }}$ row respectively.
Table 3.2: Maximum Row Table (MRT )

|  | D1 | D2 | D3 | Sources |
| :---: | :---: | :---: | :---: | :---: |
| S1 | $8-6=2$ | $8-8=0$ | $8-4=4$ | 14 |
| S2 | $9-4=5$ | $9-9=0$ | $9-8=1$ | 12 |
| S3 | $6-1=5$ | $6-2=4$ | $6-6=0$ | 5 |
| Destinations | 6 | 10 | 15 | 31 |

For MCT, subtract $6,9,8$ from each elements of $1^{\text {st }}, 2^{\text {nd }}$ and $3^{\text {rd }}$ column respectively.
Table 3.3: Maximum Column Table (MCT )

|  | D1 | D2 | D3 | Sources |
| :---: | :---: | :---: | :---: | :---: |
| S1 | $6-6=0$ | $9-8=1$ | $8-4=4$ | 14 |
| S2 | $6-4=2$ | $9-9=0$ | $8-8=0$ | 12 |
| S3 | $6-1=5$ | $9-2=7$ | $8-6=2$ | 5 |
| Destinations | 6 | 10 | 15 | 31 |

Step 2: Add maximum row table and maximum column table and name the new table as revised table.

Table 3.4: Revised Table $=$ MRT + MCT

|  | D1 | D2 | D3 | Sources |
| :---: | :---: | :---: | :---: | :---: |
| S1 | $2+0=\mathbf{2}$ | $0+1=\mathbf{1}$ | $4+4=\mathbf{8}$ | 14 |
| S2 | $5+2=\mathbf{7}$ | $0+0=\mathbf{0}$ | $1+0=\mathbf{1}$ | 12 |
| S3 | $5+5=\mathbf{1 0}$ | $4+7=\mathbf{1 1}$ | $0+2=\mathbf{2}$ | 5 |
| Destinations | 6 | 10 | 15 | 31 |

Step 3: Allocations in revised table are shown in table 3.5
Table 3.5: Allocations in Revised Table
$\left.\begin{array}{|c|c|c|c|c|}\hline & \text { D1 } & \text { D2 } & \text { D3 } & \text { Supply } \\ \hline \text { S1 } & 2 & 1 & 8 & \mathbf{\mathbf { 4 }}\end{array}\right] 14$

- First Allocation: In table 3.1.5, we see that 11 is the largest cost so we allocate 5 in cell $(3,2)$ of supply and cross-out third row. The demand of D2 is reduced to $5=(10-$ 5)
- Second Allocation: In the table next largest cost is 8 so we allocate 14 in cell $(1,3)$ of supply and cross-out first row. The demand of D3 is reduced to $1=(15-14)$
- By considering the process, third allocation in cell $(2,1)$ is 6 , fourth allocation in cell $(2,3)$ is 1 and fifth allocation in cell $(2,2)$ is 5

Step 4 : Make the same positioned allotment in the original table as revised table.
Table 3.8: Allocations in original table.
$\left.\begin{array}{|c|c|c|c|c|}\hline & \text { D1 } & \text { D2 } & \text { D3 } & \text { Sources } \\ \hline \text { S1 } & 6 & 8 & { }^{4} & \mathbf{\mathbf { 4 }}\end{array}\right] 14$

$$
Z=4 \times 6+9 \times 5+2 \times 5+4 \times 14+8 \times 1=143
$$

## 4. Optimality test for example using Modi method

## $1{ }^{\text {stIIteration }}$

Step: 1 Using table 3.8. Initial Basic Feasible Solution
Intend the values of dual variables $U_{i}$ and $V_{j}, U_{i}+V_{j}=C_{i j}$
The initially we take $\mathrm{U}_{1}=0$

| $\mathrm{C}_{13}=\mathrm{U}_{1}+\mathrm{V}_{3}=4$ | $\Rightarrow$ | $0+\mathrm{V}_{3}=4$ | $\Rightarrow$ | $\mathrm{~V}_{3}=4$ |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{C}_{21}=\mathrm{U}_{2}+\mathrm{V}_{1}=4$ | $\Rightarrow$ | $4+\mathrm{V}_{1}=4$ | $\Rightarrow$ | $\mathrm{~V}_{1}=0$ |
| $\mathrm{C}_{22}=\mathrm{U}_{2}+\mathrm{V}_{2}=9$ | $\Rightarrow$ | $4+\mathrm{V}_{2}=9$ | $\Rightarrow$ | $\mathrm{~V}_{2}=5$ |
| $\mathrm{C}_{23}=\mathrm{U}_{2}+\mathrm{V}_{3}=8$ | $\Rightarrow$ | $\mathrm{U}_{2}+4=8$ | $\Rightarrow$ | $\mathrm{U}_{2}=4$ |
| $\mathrm{C}_{32}=\mathrm{U}_{3}+\mathrm{V}_{2}=2$ | $\Rightarrow$ | $\mathrm{U}_{3}+5=2$ | $\Rightarrow$ | $\mathrm{U}_{3}=-3$ |

Table 4.1 : Example-1: Step 1 of $1^{\text {st }}$ iteration

|  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | Supply | $\mathrm{U}_{\mathrm{i}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{S}_{1}$ | 6 | 8 | 4 | 14 | $\mathrm{U}_{1}=0$ |
| $\mathrm{~S}_{2}$ | 4 | 9 | 8 | 12 | $\mathrm{U}_{2}=4$ |
| $\mathrm{~S}_{3}$ | 1 | 2 | 6 | 5 | $\mathrm{U}_{3}=-3$ |
| Demand | 6 | 10 | 15 | $\sum 31$ |  |
| $\mathrm{~V}_{\mathrm{j}}$ | $\mathrm{V}_{1}=0$ | $\mathrm{~V}_{2}=5$ | $\mathrm{~V}_{3}=4$ |  |  |

Count the occasion test using Modi method

| $\mathrm{T}_{11}=\mathrm{C}_{11}-\left(\mathrm{U}_{1}+\mathrm{V}_{1}\right)$ | $\Rightarrow$ | $6-(0+0)$ | $\Rightarrow$ | $\mathrm{T}_{11}=6$ |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{~T}_{12}=\mathrm{C}_{12}-\left(\mathrm{U}_{1}+\mathrm{V}_{2}\right)$ | $\Rightarrow$ | $8-(0+5)$ | $\Rightarrow$ | $\mathrm{T}_{12}=3$ |
| $\mathrm{~T}_{31}=\mathrm{C}_{31}-\left(\mathrm{U}_{3}+\mathrm{V}_{1}\right)$ | $\Rightarrow$ | $1-(-3+0)$ | $\Rightarrow$ | $\mathrm{T}_{31}=4$ |
| $\mathrm{~T}_{33}=\mathrm{C}_{32}-\left(\mathrm{U}_{3}+\mathrm{V}_{3}\right)$ | $\Rightarrow$ | $6-(-3+4)$ | $\Rightarrow$ | $\mathrm{T}_{33}=5$ |

Test the all placed cell sign
Table 4.2 : Example-1: Step 2 of $1^{\text {st }}$ iteration

|  | $\mathrm{D}_{1}$ | $\mathrm{D}_{2}$ | $\mathrm{D}_{3}$ | Supply | $\mathrm{U}_{\mathrm{i}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\mathrm{S}_{1}$ | 6 <br> 6 | 8 <br> 3 | $\begin{aligned} & 4 \\ & \mathbf{1 4} \end{aligned}$ | 14 | $\mathrm{U}_{1}=0$ |
| S 2 | 4 $\underline{6}$ | 9 <br> 5 | $\begin{aligned} & 8 \\ & \underline{1} \end{aligned}$ | 12 | $\mathrm{U}_{2}=4$ |
| $\mathrm{S}_{3}$ | 4 | 2 <br> 5 | $\begin{array}{r} 6 \\ 5 \end{array}$ | 5 | $\mathrm{U}_{3}=-3$ |
| Demand | 6 | 10 | 15 | $\sum 31$ |  |
| $\mathrm{V}_{\mathrm{j}}$ | $\mathrm{V}_{1}=0$ | $\mathrm{V}_{2}=5$ | $\mathrm{V}_{3}=4$ |  |  |

All unoccupied cell sign are positive or equal to zero $\mathrm{T} \geq 0$

Table 4.3: Example-1: Step 3 of $1^{\text {st }}$ iteration

|  | D1 | D2 | D3 | Sources |
| :---: | :---: | :---: | :---: | :---: |
| S1 | 6 | 8 | 4 | $\mathbf{\mathbf { 1 4 }}$ |
| S2 | 4 | $\mathbf{6}$ | 9 | $\mathbf{5}$ |

$$
Z=4 \times 14+4 \times 6+9 \times 5+8 \times 1+2 \times 5=143
$$

## 5. Results and Discussion

We have investigated the performance of proposed method by comparing its results with NWCM, LCM and VAM by testing six transportation problems, the results obtained using proposed in examples 2, 3, 5 and 6 are same as the optimal outcome. While in example 1 and 4 it is close to the optimal solution. It could be clearly seen that the proposed method gave reliable results as compared to the well-known existing methods like NWCM, LCM and VAM.

| No. of <br> Example | Type of <br> problem | Result of <br> NWCM | Result of <br> LCM | Result of <br> VAM | Result of <br> proposed <br> method | Optimal <br> Solution |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $3 \times 4$ | 116 | 116 | 101 | 101 | 100 |
| 2 | $3 \times 3$ | 39 | 37 | 33 | 33 | 33 |
| 3 | $3 \times 3$ | 730 | 555 | 555 | 555 | 555 |
| 4 | $3 \times 4$ | 3680 | 3680 | 3520 | 3510 | 3460 |
| 5 | $3 \times 3$ | 163 | 228 | 143 | 143 | 143 |
| 6 | $3 \times 3$ | 6600 | 6540 | 5920 | 5920 | 5920 |

## 6. Conclusion

In this paper, for achieving better initial basic feasible solution of transportation problems, an effective approach of Four-Step Method for optimal solution of transportation problem has been developed. The proposed method has been checked on various transportation problems and also looked over for optimality. A compression of proposed method has been made with LCM, NWCM and VAM by examining six examples, it is resulted that the proposed method is achieving better outcomes as compared to the well-known existing methods mentioned above.

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